



ON MAXIMIZATION OF PROFITS IN SOME ESTABLISHMENTS USING ANT COLONY OPTIMIZATION

¹Oke, M.O., ²Elegbede, O.O., ³Raji, R.A., ⁴Adenipekun, A.E. and ⁵Adetolaju, O.S.

^{1,2}Department of Mathematics, Ekiti State University, Ado-Ekiti, Nigeria

³Department of Mathematics & Statistics, Osun State Polytechnic, Iree, Nigeria

⁴Department of Statistics, Federal Polytechnic Ede, Osun State, Nigeria

⁵Department of Computer Science, Ekiti State University, Ado-Ekiti, Nigeria

ABSTRACT:

Many problems of combinatorial optimization arise in engineering, science, management and social sciences in which a set of optimal solution is required for a function represented as a tuple $\rho = (S, F)$. Ant colony optimization is an heuristics optimization technique for solving combinatorial optimization problems. In this paper, a metaheuristic algorithm is given to solve unconstrained optimization problems involving maximization of profits in some establishments. The algorithm was tested on four classes of problems from different establishments with paths ranging from 9 to 41 and ants ranging from 4 to 20. The superiority of Ant colony optimization over other methods was shown when we compared our results with the solutions obtained using Fibonacci search method as it performs better in all the problems considered.

KEYWORDS: Ant Colony Optimization, Chemical Pheromone Trails, Combinatorial Optimization, Fibonacci search method, Swarm Intelligent, Turple.

1.0 INTRODUCTION

Optimization is the act of obtaining the best result under any given circumstances. Engineers and scientists normally take many technological and managerial decisions at several stages of design, construction, maintenance and managerial systems. The ultimate goals of such decisions are to either minimize the efforts required or maximize the desired benefits, [11, 12]. Since the efforts required or the benefits desired in any practical situation can be expressed as functions of certain decision variables, optimization can be defined as the process of finding the conditions that give the maximum or minimum value of a function, [12].

Modern optimization techniques, sometimes called nontraditional optimization methods, have emerged as powerful and popular methods for solving complex optimization problems in recent years, [9, 10]. Ant colony optimization (ACO) is a technique of optimization that was introduced in the early 1990's by Dorigo for solving combinatorial optimization problems. The inspiring source of ACO is the foraging behaviour of real ant colonies. This behaviour is exploited in

artificial ant colonies for the search of approximate solutions to discrete and continuous optimization problems, [3, 4, 6]. The agents, called ants, are very efficient at sampling the problem space and quickly finding good solutions to it. Ant algorithms are multi-agent systems in which the behaviour of each single agent, called artificial ant, is inspired by the behaviour of real ants, [1, 3, 5]. At the core of this behaviour is the indirect communication between the ants by means of chemical pheromone trails, which enable them to find short paths between their nest and food sources, [3, 4]. This characteristic of real ant colonies is exploited in ACO algorithms in order to solve optimization problems, [4, 7].

A combinatorial optimization problem can be represented as a tuple $\rho = (S, F)$, where S is the solution space with $s \in S$ been a specific candidate solution and $F: S \rightarrow \mathbb{R}_+$ a fitness function assigning strictly positive values to candidate solutions, [3, 13]. In this case, higher values correspond to better solutions $S^* \subseteq S$ that maximize the fitness function. The solution s^* is then called an optimal solution and S^* is called the set of optimal solutions.

A lot of researches had been carried out on this modern or nontraditional method of optimization. Stützle and Hoos [14] developed an algorithm which gave some improvements on ant system while Stützle and Dorigo [13] gave some convergence proofs for a class of ant colony optimization algorithms. Dorigo and Gambardella [4], Gambardella and Dorigo [7] and Lawler *et al.* [8] applied the modern methods of optimization to solve the traveling salesman problems. Dorigo *et al.* [5], Oke [9] and Oke [10] applied some modern optimization techniques to solve complex engineering problems.

We are therefore motivated by the fact that none of these researchers have worked on the application of ACO algorithms in maximizing profits. Therefore, in this research, we considered the maximization of profits in some establishments using ant colony optimization techniques.

2.0 MATERIALS AND METHODS

In applying the ACO algorithm to unconstrained optimization problems for the maximization of profits in some establishments, we considered a profit function of the form:

Profit Function = Revenue Function - Cost Function, [2]

That is

$$P(X) = R(X) - C(X) \tag{1}$$

where $P(X)$ = Profit function

$R(X)$ = Revenue function

$C(X)$ = Cost function, [2]

In the maximization of profits in the establishment considered, the following steps were followed in applying the ACO algorithm for the solution of the problems:

Step 1: Assume a suitable number of ants in the colony (N) and a set of permissible discrete values for each of the n design variables. Denote the permissible discrete values of the design variable x_i as x_1, x_2, \dots, x_{ip} ($i = 1, 2, \dots, n$). Assume equal amounts of pheromone $\tau_{ij}^{(0)}$ initially along all the arcs or rays (discrete values of design variables) of the multilayered graph. The superscript to τ_{ij} denote the iteration number and for simplicity, let $\tau_{ij}^{(0)} = 1$ be assumed for all arcs ij and set the iteration number $l = 1$.

Step 2: (a) Compute the probability (p_{ij}) of selecting the arc or ray (or the discrete value) x_{ij} as

$$p_{ij} = \frac{\tau_{ij}^{(l)}}{\sum_{m=1}^p \tau_{im}^{(l)}}, i = 1, 2, 3, \dots, n; j = 1, 2, 3, \dots, p \quad (2)$$

(b) The specific path (or discrete values) chosen by the k th ant can be determined using random numbers generated in the range (0, 1). For this, we find the cumulative probability ranges associated with different paths based on the probabilities given by equation (2). The specific path chosen by ant k will be determined using the roulette-wheel selection process in step 3(a).

Step 3: (a) Generate N random numbers r_1, r_2, \dots, r_N in the range (0, 1), one for each ant. Determine the discrete value or path assumed by ant k for variable I as the one for which the cumulative probability range [found in step 2(b)] includes the value r_i .

(b) Repeat step 3(a) for all design variables $i = 1, 2, \dots, n$.

(c) Evaluate the objective function values corresponding to the complete paths (design vectors $X^{(k)}$ or values of x_{ij} chosen for all design variables $i = 1, 2, \dots, n$ by ant k , $k = 1, 2, \dots, N$):

$$f_k = f(X^{(k)}); k = 1, 2, \dots, N \quad (3)$$

Determine the best and worst paths among the N paths chosen by different ants:

$$f_{best} = \min_{k=1, 2, \dots, N} [f_k] \quad (4)$$

$$f_{worst} = \max_{k=1,2,\dots,N} [f_k] \quad (5)$$

Step 4: Test for the convergence of the process. The process is assumed to have converged if all N ants take the same best path. If convergence is not achieved, assume that all the ants return home and start again in search of food. Set the new iteration number as $l = l + 1$, and update the pheromones on different arcs (or discrete values of design variables) as

$$\tau_{ij}^{(l)} = \tau_{ij}^{(old)} + \sum_k \Delta\tau_{ij}^{(k)} \quad (6)$$

where $\tau_{ij}^{(old)}$ denotes the pheromone amount of the previous iteration left after evaporation, which is taken as

$$\tau_{ij}^{(old)} = (1 - \rho)\tau_{ij}^{(l-1)} \quad (7)$$

and $\Delta\tau_{ij}^{(k)}$ is the pheromone deposited by the best ant k on its path and the summation extends over all the best ants k (if multiple ants take the same best path). Note that the best path involves only one arc i, j (out of p possible arcs) for the design variable i . The evaporation rate or pheromone decay factor ρ is assumed to be in the range 0.5 to 0.8 and the pheromone deposited $\Delta\tau_{ij}^{(k)}$ is computed using the equation

$$\Delta\tau_{ij}^{(k)} = \begin{cases} \zeta f_{best} & ; \text{if } (i, j) \in \text{global best tour} \\ f_{worst} & \\ 0 & ; \text{otherwise} \end{cases} \quad (8)$$

with the new values of $\tau_{ij}^{(l)}$, go to step 2. Steps 2, 3, and 4 are repeated until the process converges, that is, until all the ants choose the same best path, [3 – 6].

3.0 COMPUTATIONAL EXAMPLES

Problem 1: Maximize $f(x)$ in the range $220 \leq x \leq 300$

where $f(x)$ is the profit function of a firm given by:

$$f(x) = -200x^2 + 92000x - 8400000$$

Problem 2: Maximize $f(x)$ in the range $4700 \leq x \leq 6500$

where $f(x)$ is the profit function of a company given by:

$$f(x) = -0.000004x^3 + 0.025x^2 + 100x - 75000$$

Problem 3: Maximize $f(x)$ in the range $150 \leq x \leq 250$

where $f(x)$ is the profit function of a firm given by:

$$f(x) = -8x^2 + 3200x - 80000$$

Problem 4: Maximize $f(x)$ in the range $0 \leq x \leq 80$

where $f(x)$ is the profit function of an industry given by:

$$f(x) = -2x^2 + 200x - 2000$$

4.0 RESULTS AND DISCUSSIONS

The solution to the problems using Ant Colony Optimization technique and the comparison between Ant Colony Optimization and Fibonacci Search methods were presented in this section.

Table 1: Ant Colony Optimization Method for Problem 1

$Ant = 4, n = 1, x = x_1$, is assumed within the range of x_1 as $(p = 9) x_{ij}; j = 1, 2, \dots, 9$

$$f(x) = -200x^2 + 92000x - 8400000, \text{ in the range } 220 \leq x \leq 300$$

| Iterations | Number of ants on best path | x_{best} | x_{worst} | f_{best} | f_{worst} |
|------------|-----------------------------|----------------|----------------|------------|-------------|
| 1 | 1 | $x_{13} = 240$ | $x_{18} = 290$ | 2160000 | 1460000 |
| 2 | 1 | $x_{12} = 230$ | $x_{17} = 280$ | 2180000 | 1680000 |
| 3 | 3 | $x_{12} = 230$ | $x_{17} = 280$ | 2180000 | 1680000 |
| 4 | 4 | $x_{12} = 230$ | | 2180000 | |

Table 2: Fibonacci Search Method for Problem 1

$$a = x_1 = 300, \quad b = x_3 = 220, \quad n = 20, \quad \epsilon = 0$$

$$f(x) = -200x^2 + 92000x - 8400000, \text{ in the range } 220 \leq x \leq 300$$

| Itera- | f_n | f_{n-1} | x_1 | x_2 | x_3 | x_4 | $f(x_2)$ | $f(x_4)$ |
|--------|-------|-----------|-------|-------|-------|-------|----------|----------|
|--------|-------|-----------|-------|-------|-------|-------|----------|----------|

| tions | | | | | | | | |
|-------|-------|------|-------------|------------|------------|------------|-----------|------------|
| 1 | 10946 | 6765 | 300 | 220 | 250.557281 | 269.442719 | 2095479.7 | 1868854.39 |
| 2 | 6765 | 4181 | 250.5572812 | 220 | 231.671843 | 238.885438 | 2179441.0 | 2164209.8 |
| 3 | 4181 | 2584 | 231.6718425 | 220 | 224.458247 | 227.213595 | 2173859.8 | 2178447.19 |
| 4 | 2584 | 1597 | 231.6718425 | 227.213595 | 228.916494 | 229.968944 | 2179765.2 | 2179999.81 |
| 5 | 1597 | 987 | 231.6718425 | 229.968944 | 230.619394 | 231.021393 | 2179923.3 | 2179791.35 |
| 6 | 987 | 610 | 230.6193935 | 229.968944 | 230.217950 | 230.370988 | 2179990.6 | 2179972.47 |
| 7 | 610 | 377 | 230.2173497 | 229.968944 | 230.063827 | 230.122467 | 2179999.2 | 2179997 |
| 8 | 377 | 233 | 230.0638267 | 229.968944 | 230.005186 | 230.027585 | 2180000 | 2179999.85 |
| 9 | 233 | 144 | 230.0051855 | 229.968944 | 229.982787 | 229.991342 | 2179999.9 | 2179999.98 |
| 10 | 144 | 89 | 230.0051855 | 229.991342 | 229.996630 | 229.999898 | 2179999.9 | 2180000 |
| 11 | 89 | 55 | 230.0051855 | 229.999898 | 230.001918 | 230.003166 | 2180000 | 2180000 |

Table 3: Ant Colony Optimization Method for Problem 2

$Ant = 20, n = 1, x = x_1$, is assumed within the range of x_1 as $(p = 41) x_{ij}; j = 1, 2, \dots, 41$

$f(x) = -0.000004x^3 + 0.025x^2 + 100x - 75000$, in the range $4700 \leq x \leq 6500$

| Iterations | Number of ants on best path | x_{best} | x_{worst} | f_{best} | f_{worst} |
|------------|-----------------------------|-------------------|-------------------|-------------|-------------|
| 1 | 1 | $x_{1,22} = 5645$ | $x_{1,1} = 4700$ | 566615.7805 | 531958 |
| 2 | 1 | $x_{1,22} = 5645$ | $x_{1,41} = 6500$ | 566615.7805 | 532750 |
| 3 | 6 | $x_{1,22} = 5645$ | $x_{1,40} = 6455$ | 566615.7805 | 536333.04 |
| 4 | 17 | $x_{1,22} = 5645$ | $x_{1,5} = 4880$ | 566615.7805 | 543502.912 |
| 5 | 20 | $x_{1,22} = 5645$ | | 566615.7805 | |

Table 4: Fibonacci Search Method for Problem 2

$a = x_1 = 6500, b = x_3 = 4700, n = 20, \epsilon = 0$

$f(x) = -0.000004x^3 + 0.025x^2 + 100x - 75000$, in the range $4700 \leq x \leq 6500$

| Iterations | f_n | f_{n-1} | x_1 | x_2 | x_3 | x_4 | $f(x_2)$ | $f(x_4)$ |
|------------|-------|-----------|------------|------------|------------|------------|------------|------------|
| 1 | 10946 | 6765 | 6500 | 4700 | 5387.53883 | 5812.46117 | 563887.607 | 565374.598 |
| 2 | 6765 | 4181 | 6500 | 5812.46117 | 6075.07763 | 6237.38354 | 558330.901 | 550701.725 |
| 3 | 4181 | 2584 | 6075.07763 | 5812.46117 | 5912.77174 | 5974.76706 | 563436.347 | 561777.593 |
| 4 | 2584 | 1597 | 5912.77174 | 5812.46117 | 5850.77639 | 5874.45652 | 564741.866 | 564284.531 |
| 5 | 1597 | 987 | 5850.77693 | 5812.46117 | 5827.09629 | 5836.14128 | 565148.497 | 564999.142 |
| 6 | 987 | 610 | 5827.09629 | 5812.46117 | 5818.05213 | 5821.50618 | 565290.502 | 565237.128 |
| 7 | 610 | 377 | 5818.05128 | 5812.46117 | 5814.59641 | 5815.91605 | 565342.806 | 565322.954 |

| | | | | | | | | |
|----|-----|-----|------------|------------|------------|------------|------------|------------|
| 8 | 377 | 233 | 5814.59641 | 5812.46117 | 5813.27676 | 5813.78083 | 565362.503 | 565354.998 |
| 9 | 233 | 144 | 5813.27676 | 5812.46117 | 5812.77270 | 5812.96522 | 565369.985 | 565367.130 |
| 10 | 144 | 89 | 5812.77270 | 5812.46117 | 5812.58016 | 5812.73227 | 565372.837 | 565370.584 |
| 11 | 89 | 55 | 5812.58016 | 5812.46117 | 5812.50663 | 5812.53471 | 565373.925 | 565373.510 |
| 12 | 55 | 34 | 5812.50663 | 5812.46117 | 5812.47853 | 5812.48927 | 565374.341 | 565374.182 |
| 13 | 34 | 21 | 5812.47853 | 5812.46117 | 5812.46781 | 5812.47189 | 565374.500 | 565374.439 |
| 14 | 21 | 13 | 5812.46781 | 5812.46117 | 5812.46370 | 5812.47192 | 565374.561 | 565374.438 |
| 15 | 13 | 8 | 5812.46370 | 5812.46117 | 5812.46215 | 5812.46273 | 565374.584 | 565374.575 |
| 16 | 8 | 5 | 5812.46215 | 5812.46117 | 5812.46154 | 5812.46178 | 565374.593 | 565374.589 |
| 17 | 5 | 3 | 5812.46154 | 5812.46117 | 5812.46132 | 5812.46139 | 565374.596 | 565374.595 |
| 18 | 3 | 2 | 5812.46132 | 5812.46117 | 5812.46122 | 5812.46127 | 565374.597 | 565374.597 |
| 19 | 2 | 1 | 5812.46122 | 5812.46117 | 5812.46120 | 5812.46120 | 565374.598 | 565374.598 |

Table 5: Ant Colony Optimization Method for Problem 3

$Ant = 8, n = 1, x = x_1$, is assumed within the range of x_1 as $(p = 11) x_{ij}; j = 1, 2, \dots, 11$

$f(x) = -8x^2 + 3200x - 80000$, in the range $150 \leq x \leq 250$

| Iterations | Number of ants on best path | x_{best} | x_{worst} | f_{best} | f_{worst} |
|------------|-----------------------------|-----------------|--------------------------------|------------|-------------|
| 1 | 1 | $x_{1,5} = 200$ | $x_{1,1}, x_{1,11} = 150, 250$ | 240000 | 220000 |
| 2 | 3 | $x_{1,5} = 200$ | $x_{1,1} = 150$ | 240000 | 220000 |
| 3 | 6 | $x_{1,5} = 200$ | $x_{1,9} = 230$ | 240000 | 232800 |
| 4 | 7 | $x_{1,5} = 200$ | $x_{1,2} = 160$ | 240000 | 227200 |
| 5 | 8 | $x_{1,5} = 200$ | | 240000 | |

Table 6: Fibonacci Search Method for Problem 3

$a = x_1 = 250, b = x_3 = 140, n = 20, \epsilon = 0$

$f(x) = -8x^2 + 3200x - 80000$, in the range $150 \leq x \leq 250$

| Iterations | c_n | f_{n-1} | x_1 | x_2 | x_3 | x_4 | $f(x_2)$ | $f(x_4)$ |
|------------|-------|-----------|------------|------------|------------|------------|------------|------------|
| 1 | 10946 | 6765 | 250 | 140 | 182.016262 | 207.983738 | 237412.681 | 239490.079 |
| 2 | 6765 | 4181 | 250 | 207.983738 | 224.032522 | 233.951217 | 235379.503 | 230778.519 |
| 3 | 4181 | 2584 | 224.032522 | 207.983738 | 214.113829 | 217.902432 | 238406.399 | 237436.024 |
| 4 | 2584 | 1597 | 214.113829 | 207.983738 | 210.325224 | 211.772343 | 239147.118 | 238891.296 |
| 5 | 1597 | 987 | 210.325224 | 207.983738 | 208.878107 | 209.430856 | 239369.434 | 239288.472 |

| | | | | | | | | |
|----|-----|-----|------------|------------|------------|------------|------------|------------|
| 6 | 987 | 610 | 208.878107 | 207.983738 | 208.325356 | 208.536489 | 239445.508 | 239417.027 |
| 7 | 610 | 377 | 208.325356 | 207.983738 | 208.114225 | 208.194870 | 239473.275 | 239462.753 |
| 8 | 377 | 233 | 208.114225 | 207.983738 | 208.033580 | 208.064384 | 239483.693 | 239479.726 |
| 9 | 233 | 144 | 208.033580 | 207.983738 | 208.002776 | 208.014542 | 239487.645 | 239486.137 |
| 10 | 144 | 89 | 208.002776 | 207.983738 | 207.991010 | 207.995505 | 239489.150 | 239488.575 |
| 11 | 89 | 55 | 207.991010 | 207.983738 | 207.986516 | 207.988232 | 239489.725 | 239489.505 |
| 12 | 55 | 34 | 207.986516 | 207.983738 | 207.984799 | 207.985456 | 239489.944 | 239489.86 |
| 13 | 34 | 21 | 207.984799 | 207.983738 | 207.984144 | 207.984394 | 239490.028 | 239489.996 |
| 14 | 21 | 13 | 207.984144 | 207.983738 | 207.983893 | 207.983990 | 239490.060 | 239490.047 |
| 15 | 13 | 8 | 207.983893 | 207.983738 | 207.983798 | 207.983834 | 239490.072 | 239490.067 |
| 16 | 8 | 5 | 207.983798 | 207.983738 | 207.983761 | 207.983776 | 239490.077 | 239490.075 |
| 17 | 5 | 3 | 207.983761 | 207.983738 | 207.983747 | 207.983752 | 239490.078 | 239490.078 |
| 18 | 3 | 2 | 207.983747 | 207.983738 | 207.983741 | 207.983744 | 239490.079 | 239490.079 |

Table 7: Ant Colony Optimization Method for Problem 4

$Ant = 8, n = 1, x = x_1$, is assumed within the range of x_1 as $(p = 17) x_{ij}$;

$j = 1, 2, \dots, 17$

$f(x) = -2x^2 + 200x - 2000$, in the range $0 \leq x \leq 80$

| Iterations | Number of ants on best path | x_{best} | x_{worst} | f_{best} | f_{worst} |
|------------|-----------------------------|-----------------|-------------------------------|------------|-------------|
| 1 | 1 | $x_{1,11} = 50$ | $x_{14} = 15$ | 3000 | 550 |
| 2 | 14 | $x_{1,11} = 50$ | $x_{14} = 15$ | 3000 | 550 |
| 3 | 18 | $x_{1,11} = 50$ | $x_{1,12}, x_{1,10} = 55, 45$ | 3000 | 2950 |
| 4 | 20 | $x_{1,11} = 50$ | | 3000 | |

Table 8: Fibonacci Search Method for Problem 4

$a = x_1 = 80, b = x_3 = 0, n = 20, \epsilon = 0$

$f(x) = -2x^2 + 200x - 2000$, in the range $0 \leq x \leq 80$

| Iterations | f_n | f_{n-1} | x_1 | x_2 | x_2 | x_4 | $f(x_2)$ | $f(x_4)$ |
|------------|-------|-----------|-----------|----------|------------|------------|------------|-------------|
| 1 | 10946 | 6765 | 80 | 0 | 30.5572812 | 49.4427188 | 2243.96137 | 2999.378875 |
| 2 | 6765 | 4181 | 80 | 49.44272 | 61.1145132 | 68.3281575 | 2752.93305 | 2328.157291 |
| 3 | 4181 | 2584 | 61.114561 | 49.44272 | 53.9009662 | 56.6563139 | 2969.56493 | 2969.564929 |
| 4 | 2584 | 1597 | 53.900966 | 49.44272 | 51.1456175 | 52.1980675 | 2997.37512 | 2990.337 |
| 5 | 1597 | 987 | 51.145618 | 49.44272 | 50.0931682 | 50.4951678 | 2999.98264 | 2999.509614 |

| | | | | | | | | |
|----|-----|-----|-----------|----------|------------|------------|------------|-------------|
| 6 | 987 | 610 | 50.093169 | 49.44272 | 49.6911682 | 49.8447191 | 2999.80925 | 2999.951776 |
| 7 | 610 | 377 | 50.093169 | 49.84472 | 49.9396187 | 49.998269 | 2999.99271 | 2999.999994 |
| 8 | 377 | 233 | 50.093169 | 49.99827 | 50.0345171 | 50.0569204 | 2999.99762 | 2999.993524 |
| 9 | 233 | 144 | 50.034517 | 49.99827 | 50.0121149 | 50.0206713 | 2999.99971 | 2999.999145 |
| 10 | 144 | 89 | 50.012115 | 49.99827 | 50.0035574 | 50.0068265 | 2999.99998 | 2999.999907 |
| 11 | 89 | 55 | 50.003557 | 49.99827 | 50.0002893 | 50.0015371 | 2999.99999 | 2999.999999 |
| 12 | 55 | 34 | 50.003557 | 50.00154 | 50.0023085 | 50.0027860 | 2999.99998 | 2999.999998 |
| 13 | 34 | 21 | 50.002309 | 50.00154 | 50.0018320 | 50.0020135 | 2999.99998 | 2999.999999 |
| 14 | 21 | 13 | 50.002309 | 50.00201 | 50.0021259 | 50.0021961 | 2999.99998 | 2999.999992 |
| 15 | 13 | 8 | 50.002309 | 50.00220 | 50.0022393 | 50.0022652 | 2999.99998 | 2999.999994 |
| 16 | 8 | 5 | 50.002309 | 50.00227 | 50.0022814 | 50.0022922 | 2999.99999 | 2999.999991 |
| 17 | 5 | 3 | 50.002281 | 50.00227 | 50.0022717 | 50.0022275 | 2999.99998 | 2999.999988 |

Table 9: Comparison between ACO and Fibonacci Search Method for Problem 1

| | ACO method | Fibonacci method |
|------------|------------------------------|------------------------------|
| Function | $-200p^2 + 92000p - 8400000$ | $-200p^2 + 92000p - 8400000$ |
| Range | $220 \leq x \leq 300$ | $220 \leq x \leq 300$ |
| x_{best} | 230 | 230.001918, 230.0031656 |
| f_{best} | 2180000 | 2180000 |

Table 10: Comparison between ACO and Fibonacci Search Method for Problem 2

| | ACO method | Fibonacci method |
|------------|--|--|
| Function | $-75000 + 100x + 0.025x^2 - 0.000004x$ | $-75000 + 100x + 0.025x^2 - 0.000004x^3$ |
| Range | $4700 \leq x \leq 6500$ | $4700 \leq x \leq 6500$ |
| x_{best} | 5645 | 5812.461198, 5812.461197 |
| f_{best} | 566615.7805 | 565374.5976 |

Table 11: Comparison between ACO and Fibonacci Search Method for Problem 3

| | ACO method | Fibonacci method |
|------------|-------------------------|--------------------------|
| Function | $-8x^2 + 3200x - 80000$ | $-8x^2 + 3200x - 80000$ |
| Range | $150 \leq x \leq 250$ | $150 \leq x \leq 250$ |
| x_{best} | 200 | 207.9837414, 207.9837443 |
| f_{best} | 240000 | 239490.079 |

Table 12: Comparison between ACO and Fibonacci Search Method for Problem 4

| | ACO method | Fibonacci method |
|------------------------------|--|--|
| Function | $200x - 2x^2 - 2000$ | $200x - 2x^2 - 2000$ |
| Range | $0 \leq x \leq 80$ | $0 \leq x \leq 80$ |
| x_{best} | 50 | 50.00227171, 50.00227496 |
| f_{best} | 3000 | 2999.999988 |

5.0 CONCLUSION

The Ant Colony Optimization has been applied in the maximization of profits in some selected establishments. The results obtained showed the superiority of Ant Colony Optimization over the Fibonacci search method as it performs better in all the selected problems.

6.0 REFERENCES

- [1] Bullnheimer, B., Hartl, R.F. and Strauss, C. (1999): A New Rank-based Version of the Ant System (A Computational Study), Central European Journal of Operations Research, 7(1), 25 - 38.
- [2] Chiang, A. and Wainwright, K. (2013): Fundamental Methods of Mathematical Economics, Mc Graw-Hill Co. Inc., USA
- [3] Darigo, M. and Blum, C. (2005): Ant Colony Optimization Theory (a Survey), Theoretical Computer Science, 344, 243 – 278.
- [4] Dorigo, M. and Gambardella, L.M. (1997): Ant Colony System (A Cooperative Learning Approach to the Traveling Salesman Problem), IEEE Transaction on Evolutionary Computation, 1, 53 - 66.
- [5] Dorigo, M., Maniezzo, V. and Colorni, A. (1991): The Ant System (An Autocatalytic Optimizing Process), Technical Report *TR91-016*, Politecnico di Milano.
- [6] Dorigo, M. and Stützle, T. (2004): Ant Colony Optimization, Cambridge MA, MIT Press.
- [7] Gambardella, L.M. and Dorigo, M. (1995): Ant-Q (A Reinforcement Learning Approach to the Travelling Salesman Problem), Proceedings of the 12th International Conference on Machine Learning, California, USA, 252 - 260.
- [8] Lawler, E., Lenstra, J.K., Rinnooy, K.A. and Shmoys, D.B. (1985): The Travelling Salesman Problem, John Wiley & Sons, New York.
- [9] Oke, M.O. (2012(a)): Solving Optimal Power Flow Problems using Partial Swarm Optimization Technique, Presented at the 23rd Annual Colloquium and Congress of the Nigerian Association of Mathematical Physics.
- [10] Oke, M.O. (2012(b)): Differential Evolution Method of Solving Optimal Power Flow Problems, Presented at the 31st Annual Conference of the Nigerian Mathematical Society.
- [11] Oke, M. O. (2014): A Second Order Method for Minimizing Unconstrained Optimization Problems, Journal of Applied Mathematics and Bioinformatics, International Scientific Press, UK, 4(3), 65 - 73.
- [12] Rao, S.S. (1998): Optimization Theory and Application, Wiley Eastern Limited.
- [13] Stützle, T. and Dorigo, M. (2002): A Short Convergence Proof for a Class of ACO Algorithms, IEEE Transactions on Evolutionary Computation, 6(4), 358 - 365.

[14] Stützle, T. and Hoos, H. (1997): Improvements on the Ant System. (Introducing MAX-MIN Ant System), Proceeding of International Conference on Artificial Neural Networks and Genetic Algorithms, Springer Verlag, 245 – 249.

© GSJ